Formation Control of Nonholonomic Vehicles Under Time Delayed Communications

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Abstract—This document considers the formation control problem for a group of nonholonomic mobile robots under time delayed communications. The agents are assumed to be working under a directed and fixed communication topology. A recently developed formation control technique, based on a consensus protocol for linear systems is revisited. The stability analysis of this control technique, in the domain of the time delays, is performed using the Cluster Treatment of Characteristic Roots paradigm under the Spectral Delay Space approach. This analysis leads to an exact declaration of the stability boundaries in the domain of the delays. This linear control law is adapted to the nonlinear dynamics of the nonholonomic carts by means of feedback linearization. Simulation results show the applicability of this method to the class of agents considered here.

Note to Practitioners—The problem dealt with in this paper is the coordinated motion of a group of mobile robots. Without a leader or a central authority, the members of the team attempt to converge into a predefined formation by exchanging information about position and velocities among themselves. The particularity here is that the communication among agents is corrupted by two time delays. One delay is considered to affect the measurement of the position of the agents and the second the determination of their velocity. We present a methodology to study the effect of the delay in the stability of the formation. This technique allows an exact declaration of the delay compositions that guarantee convergence to the formation in a finite time. The methodology is applied to a group of nonholonomic differential drive mobile robots.

Index Terms—Decentralized control, delay systems, networked control systems.

I. INTRODUCTION

COOPERATION among multiple agents is a research topic which has attracted a lot of attention in the past two decades. In order to perform a cooperative task, the members of a team have to reach an agreement on a certain variable. If they want to rendezvous at some point, they must agree upon the location of that point. If they are to travel together, they must select a heading and speed. This gives birth to the consensus problem. After the original agreement protocol of Vickseck et al. [1], expanded by the work of Jadbabaie et al. [2], many authors have worked to extend the theory and applications of the consensus problem. Olfati-Saber and Murray [3] set up a formal framework to treat the agreement problem, highlighting the importance of algebraic graph theory in this topic. They also considered the case in which time delays are present in the communication channels. This pioneering work was limited to agents modeled by single integrator dynamics, and it was soon expanded to consider cases including higher order dynamics [4]–[7] and multiple time delays [8]–[10].

Some researchers have also studied the more general formation generation or flocking problem [11], [12]. Yao et al. [11] followed an approach based on artificial potentials, whereas Olfati-Saber [12] used a consensus term for velocity agreement and a distance based repulsion term to create the spacing among agents. These methods are heavily nonlinear, despite the fact that they are applied to agents with simple single or double integrator dynamics. Following a different approach, Lin et al. [7] transformed a consensus algorithm into a formation control protocol by including the desired inter-agent distances in the control protocol. Cepeda-Gomez and Olgac [13], also expanded a consensus control law [9], [10] to transform it into a formation logic. The last two protocols preserve the linearity of the control logic, and therefore allow the use of well known techniques for the stability analysis with respect to the time delays. All these works, however, have considered agents with very simple dynamics.

Since real-life mobile agents have dynamics more complicated than the single or double integrator models commonly used, some authors have expanded the field to include the problem of cooperation of nonholonomic agents. The works of Zhai et al. [14], Listman [15] and Ghommam [16] present different approaches to solve this problem. However, very few works, besides that of Dong and Farrell [17], have treated the consensus or formation control problems for nonholonomic agents under time delayed communications. To study the stability with respect to the delay, they use an approach based on the D-subdivision method, and obtain a delay margin. This result is valid only if all the eigenvalues of the Laplacian matrix generated by the communication topology are real.

In this work, we depart from a consensus protocol affected by two time delays, expand it to include a constant forcing term to create the formation and then deploy it into nonholonomic agents. More specifically, we consider Differential-Drive Mobile Robots (DDMR). The delays are assumed to affect the mea-
measurements of the position and velocity of the agents. A feedback linearization procedure allows us to deploy the original double integrator control protocol to this class of agents, and the Cluster Treatment of Characteristic Roots (CTCR) Paradigm [18], [19] allows us to find nonconservative stability bounds in the domain of the time delays.

To summarize, the contribution of the present paper is the following. Using a feedback linearization model, we adapt a recently developed formation control technique to a group of nonholonomic robots. The formation control technique is based on a consensus protocol and assumes two time delays in the sensing of positions and velocities. By using the CTCR methodology, we determine, exactly and exhaustively, the delay combinations that guarantee stability, i.e., convergence to formation in a finite time. To the best of our knowledge, this is the first time such a formation problem under time delays is solved in an exact manner for nonholonomic agents.

The rest of this paper includes the description of the consensus protocol in Section II, the stability analysis with respect to the delays, using CTCR, in Section III, the formation control algorithm in Section IV, the feedback linearization procedure in Section V and some simulation examples in Section VI. Conclusions and future work directions are presented in Section VII.

II. CONSENSUS PROTOCOL

We consider a group of $n$ agents with double integrator dynamics of the form $\dot{x}_i(t) = u_i(t)$, where $x_i(t) \in \mathbb{R}^m$ represents the position vector of the $i$th agent and $u_i(t) \in \mathbb{R}^m$ is the input vector. It is assumed that each one of the position coordinates is independent of the other components of the $x_i$ vector, and that the $u_{i,k}$ input affects only the dynamics of the $x_i,k$ position, i.e., the dynamics of all the position coordinates are decoupled. Taking advantage of this, in the rest of this section we consider only a scalar position, with the understanding that the results can be easily expanded to the other coordinates for higher dimensional cases.

We assume the agents are communicating with some of their peers. In the most general case, the communication channels would be unidirectional, and therefore a directed graph is used to model the communication structure. Each vertex represents an agent and each edge a communication channel. The adjacency matrix of this graph, $A = [a_{ik}] \in \mathbb{R}^{n \times n}$, is a binary matrix. The $a_{ik}$ element is 1 if agent $i$ receives information from agent $k$ and 0 if not. The diagonal elements $a_{ii}$ are taken as 0 too. The $i$th row sum of $A$ is called the in-degree of the vertex $i$ [20], and it is represented by $\delta_i$. The in-degree is also the number of peers from which the $i$th member receives information. They are organized in a set called the informers of agent $i$, which is denoted by $N_i$. With the number of informers of each agent, the in-degree matrix of the communication topology is created. It consists of a diagonal matrix that contains the in-degree of node $i$ in its $i$th diagonal term: $\Delta = \text{diag}(\delta_1, \delta_2, \ldots, \delta_n)$.

It is assumed that the communication process is affected by two time delays. The first delay is in the position communication channels and the second one in the velocity information exchange. The different time delays in the velocity and position channels are justified considering that different sensors are used to measure these variables, and therefore different sensing delays will be added on top of the delays inherent to the communication channels. Since we are considering that the members of the group are all alike, assuming the delays are uniform across the network and constant also makes sense. It is also assumed that each agent is able to process its own information fast enough to consider the so called self delay negligible. This means, each agent has its current information about position and velocity, but receives only delayed information from its peers. These assumptions are based on the characteristics of the physical platform currently available to perform a practical implementation of the methodology described in this paper.

The control law is defined in this form: each agent takes the delayed position information it receives from its informers, computes the average and compares its own current position to that average; this error is penalized with a proportional control gain $P \in \mathbb{R}^+$. A similar task is performed with the velocity information received, using a derivative control gain $D \in \mathbb{R}^+$. Mathematically, the control law is represented by

$$u_i = P \left( \sum_{k \in N_i} \frac{x_k(t - \tau_1)}{\delta_i} - x_i(t) \right) + D \left( \sum_{k \in N_i} \frac{\dot{x}_k(t - \tau_2)}{\delta_i} - \dot{x}_i(t) \right) \quad (1)$$

for $i = 1, 2, \ldots, n$.

Using a concatenation of the positions and velocities of the agents as the state vector $z = [x_1, x_2, \ldots, x_n, \dot{x}_1, \ldots, \dot{x}_n]^T$, the dynamics of the multi-agent system, under the control law (1), can be presented in state space as

$$z = \begin{bmatrix} I_n \otimes [0 \ 1] \ \ -P \ D \end{bmatrix} z(t) + \begin{bmatrix} C \otimes [0 \ 0] \ P \ 0 \ 0 \ 0 \ 0 \ D \end{bmatrix} z(t - \tau_1) + \begin{bmatrix} 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ 0 \ \end{bmatrix} z(t - \tau_2) \quad (2)$$

where $I_n$ is the identity matrix of size $n \times n$, $\otimes$ represents the Kronecker product operation and the matrix $C \in \mathbb{R}^{m \times n}$ is the product of the inverse of the in-degree matrix and the adjacency matrix: $C = \Delta^{-1} A$. The complexity of the dynamics in (2) increases rapidly with the number of agents in the group, quickly becoming intractable. In order to simplify the stability analysis with respect to the time delays, the methodology introduced in [21], extended to the case of directed topologies, is used here. It is based in the following lemma.

Lemma 1–Factorization Property: The characteristic equation of the system (2) can always be expressed as the product of a $\ell$ second- and $m$ fourth-order factors as

$$Q(s, P, D, \tau_1, \tau_2) = \prod_{i=1}^{\ell+m} q_i(s, P, D, \tau_1, \tau_2, \lambda_i) =$$
where \( \lambda_i \) represent the eigenvalues of the matrix \( \mathbf{C} \), which has \( \ell \) real eigenvalues, denoted by \( i = 1, 2, \ldots, \ell \) and \( m \) complex conjugate eigenvalue pairs, \( (\lambda_i, \lambda_i^*) \), \( i = \ell + 1, \ell + 2, \ldots, \ell + m \). Then, \( n = \ell + 2m \).

The proof is omitted here due to space constraints, it can be found in [10] and [13].

As it is shown in [22], if the communication topology is connected and has at least one spanning tree, the matrix \( \mathbf{C} \) has 1 as an eigenvalue with multiplicity one. This factor introduces a root at \( s = 0 \) regardless of the delays, i.e., it always presents a rigid body mode. This factor is known to dictate the dynamics of a weighted average of the states of the agents, which we call the weighted centroid. It defines the value upon which the agents would agree if consensus is reached. Without loss of generality, this eigenvalue is denoted as \( \lambda_1 \). The other \( \ell + m - 1 \) factors in (3) represent the disagreement dynamics: if they are all stable, the consensus is reached. These results are well known in the literature, and are stated here as lemmas for completeness. The proofs of these lemmas can be found in [13].

Lemma 2–Group Behavior: Assume the communication topology has at least one spanning tree. Then, the agents in the group reach a consensus if and only if the factor (4) is stable, where

\[
q_1(s) = s^2 + Ds + P - (Dse^{-\tau_2s} + Pe^{-\tau_1s})
\]

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Lemma 3–Topologies Without Spanning Trees: If the given communication topology does not have a spanning tree, the control policy described by (1) cannot result in consensus.

III. STABILITY ANALYSIS USING CTCR

The results of the previous section show that agents reach consensus if all factors in (3) are stable. Section IV shows that this is also a necessary condition to ensure the formation stability. The general form of the factors is

\[
\left( \prod_{i=1}^{\ell} \left[ s^2 + Ds + P - \lambda_i (Dse^{-\tau_2s} + Pe^{-\tau_1s}) \right] \right) \left( \prod_{i=\ell+1}^{\ell+m} \left[ s^2 + 2Ds + (D^2 + 2P) s^2 + 2DP + P^2 - 2\Re(\lambda_i) (s^2 + Ds + P) (Dse^{-\tau_2s} + Pe^{-\tau_1s}) + \lambda_i^2 \right] \right) = 0
\]

where \( \lambda_i \) represent the eigenvalues of the \( \mathbf{C} \) matrix, which has \( \ell \) real eigenvalues, denoted by \( i = 1, 2, \ldots, \ell \) and \( m \) complex conjugate eigenvalue pairs, \( (\lambda_i, \lambda_i^*) \), \( i = \ell + 1, \ell + 2, \ldots, \ell + m \). Then, \( n = \ell + 2m \).

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Lemma 3–Topologies Without Spanning Trees: If the given communication topology does not have a spanning tree, the control policy described by (1) cannot result in consensus.
The question now is how to find the kernel curves. For this, several different methods can be used, like the Kronecker sum [24], the Rekasius substitution [18], the Spectral Delay Space (SDS) based on Rekasius substitution [19] or the SDS based on half-angle tangent substitution [9], which was used in this work. Due to space limitations, we do not include details of the approach we followed here, and we refer the reader to [9] and [10], where it is explained.

IV. FORMATION CONTROL

The consensus protocol in (1) is designed to bring the state of the agents to a common value. If a position state is considered, that would imply a collision. We can transform this behavior into a formation control structure by just adding a constant forcing term to the consensus protocol in (1). In order to design this term, we must consider the state transformation that decouples the system and factorizes its characteristic equation, as stated in Lemma 1 and detailed in [10]. The matrix used in this transformation, $T$, is the one that takes $C$ to its Jordan canonical form. Under the transformation $\xi = (T^{-1} \otimes I_2) z$, the dynamics (2) is decoupled into $\ell$ second-order subsystems and $m$ fourth-order subsystems. The characteristic equations of these subsystems are the factors of the global characteristic equation as declared in (3).

The state-space representation of each subsystem, including the extra forcing term we are introducing, has the general form of

$$\dot{\xi}_i(t) = A \xi_i(t) + B_1 \xi_i(t - \tau_1) + B_{2i} \xi_i(t - \tau_2) + \varphi_i$$

where the matrices $A$, $B_1$, and $B_{2i}$ have the following form for subsystems corresponding to real eigenvalues of $C$:

$$A = \begin{bmatrix} 0 & 1 \\ -P & -D \end{bmatrix}, \quad B_{1i} = \begin{bmatrix} 0 \\ P \end{bmatrix}, \quad B_{2i} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

and the following form for those corresponding to complex conjugate eigenvalue pairs:

$$A = \left( I_2 \otimes \begin{bmatrix} 0 & 1 \\ -P & -D \end{bmatrix} \right), \quad B_{1i} = \left( J_i \otimes \begin{bmatrix} 0 \\ P \end{bmatrix} \right), \quad B_{2i} = \left( J_i \otimes \begin{bmatrix} 0 \\ 0 \end{bmatrix} \right)$$

with $J_i$ being

$$J_i = \begin{bmatrix} \Re(\lambda_i) \\ \Im(\lambda_i) \end{bmatrix} \begin{bmatrix} \Re(\lambda_i^*) \\ \Im(\lambda_i^*) \end{bmatrix}$$

The forcing terms in (8) have the form $\varphi_i = [0, \phi_{i1}]^T$ for second-order factors and $\varphi_i = [0, \phi_{i1}, 0, \phi_{i2}]^T$ for fourth-order factors. These forcing terms are applied only to the disagreement dynamics, in order to create the spacing among the agents. The centroid dynamics is left undisturbed, and it dictates where the center of the formation will be located at the end of the exercise.

In order to study the effect of the forcing terms in (8), we take the Laplace transform and apply the final value theorem, arriving to

$$\xi_{i\infty} = -\frac{\phi_i}{P(\lambda_i - 1)} \begin{bmatrix} 1 \\ 0 \end{bmatrix}$$

for second order factors and to

$$\xi_{i\infty} = -\frac{\Re(\lambda_i) - 1 - \Im(\lambda_i)}{P(\lambda_i - 1)} \begin{bmatrix} \phi_{i1} \\ \phi_{i2} \end{bmatrix}$$

Notice that if we make the forcing terms equal to zero, the steady-state disagreement is zero, leading to consensus. When the forcing terms are present, both (12a) and (12b) imply constant relative positions among the agents and zero relative velocities.

With the values of $\xi_{i\infty}$, it is possible to obtain the corresponding steady state values for the agents positions, using $z_{i\infty} = (T \otimes I_2) \xi_{i\infty}$. Of course, the formation design process does not start from $\xi_{i\infty}$, but from $z_{i\infty}$, the desired locations of the agents in the formation. Then, $\xi_{i\infty}$ is found from $\xi_{i\infty} = (T^{-1} \otimes I_2) z_{i\infty}$ and (12) is used to find the value of the forcing terms in the transformed domain, with which the forcing vector $\varphi = [0, 0, \varphi_{i1}^2, \cdots, \varphi_{i1+m}^2]^T$ is created. The actual forcing terms, applied in the nontransformed domain are found using $F = (T \otimes I_2) \varphi$.

With the formation control term added, the dynamics (2) finally becomes

$$\dot{z} = \left( I_n \otimes \begin{bmatrix} 0 & 1 \\ -P & -D \end{bmatrix} \right) z(t) + \left( C \otimes \begin{bmatrix} 0 \\ P \end{bmatrix} \right) z(t - \tau_1) + \left( C \otimes \begin{bmatrix} 0 \\ 0 \end{bmatrix} \right) z(t - \tau_2) + F.$$

As a cautionary note, we would like to mention that the algorithm to design the formation can be seen as a brute force approach, in which each agent is assigned to a position in the formation. This kind of approach can be used for small groups of robots. However, since the assignment of the positions is not automated, the design of the formation becomes more involved for swarms of several tens of agents, as more parameters have to be considered. Furthermore, our methodology is not robust against failures of one of the members of the group. In this regard, it is significantly different from more formal methods to design formations, like those presented by Belta and Kumar [25].

V. FEEDBACK LINEARIZATION

The consensus protocol of Section II and the formation control algorithm of Section IV were devised for second-order holonomic agents, in which the motion along different coordinates is independent. This would be the case of some types of omnidirectional mobile robots. In this work, however, we propose the use of nonholonomic wheeled robots, more specifically, of
differential drive mobile robots, DDMR, like the iRobot Create [26]. In order to apply the protocol of previous sections to this class of robots, some transformation must be undertaken first.

The kinematics of a DDMR is similar to that of the unicycle. The variables used for its description are the coordinates of the center and the orientation of the wheel, pictured in Fig. 1. Since the wheel is restricted to rolling without slipping, the relationship between the speed, \( v \), and the angular rate \( \omega \) of the unicycle and the derivatives of its state variables is given by

\[
\begin{bmatrix}
\dot{x} \\
\dot{y} \\
\dot{\theta}
\end{bmatrix} = 
\begin{bmatrix}
\cos \theta & 0 & v \\
\sin \theta & 0 & 0 \\
0 & 1 & \omega
\end{bmatrix}
\begin{bmatrix}
v \\
\omega
\end{bmatrix}.
\] (14)

It is well known that (14) does not satisfy Brockett’s condition [27], and therefore there is no continuous, linear time invariant feedback law that can stabilize the system. The consensus protocol (1) cannot be directly applied. However, using a feedback linearization technique, proposed by [28], we can find an equivalent linear system in which the second derivative of the state equals the control input, as was assumed in Section II.

To start the feedback linearization of (14), consider a new state vector \( \mathbf{w} = [w_1, w_2]^T \). This vector is selected such that the derivatives of its components are equal to the derivatives of \( x \) and \( y \)

\[
\begin{bmatrix}
\dot{w}_1 \\
\dot{w}_2
\end{bmatrix} = 
\begin{bmatrix}
\dot{x} \\
\dot{y}
\end{bmatrix} = 
\begin{bmatrix}
v \cos \theta \\
v \sin \theta
\end{bmatrix} = v
\begin{bmatrix}
\cos \theta \\
\sin \theta
\end{bmatrix}.
\] (15)

Taking the second derivative of \( \mathbf{w} \), we get

\[
\begin{bmatrix}
\ddot{w}_1 \\
\ddot{w}_2
\end{bmatrix} = 
\begin{bmatrix}
\dot{\theta} \cos \theta - v \omega \sin \theta \\
\dot{\theta} \sin \theta + v \omega \cos \theta
\end{bmatrix}.
\] (16)

If we make the second derivative of \( \mathbf{w} \) equal to the system’s input, we obtain the equivalent linear system

\[
\begin{bmatrix}
\ddot{w}_1 \\
\ddot{w}_2
\end{bmatrix} = 
\begin{bmatrix}
\dot{u}_1 \\
\dot{u}_2
\end{bmatrix}.
\] (17)

By using (16) and (17), we obtain a relation between the input to the linear system and that of the nonlinear system

\[
\begin{bmatrix}
u_1 \\
u_2
\end{bmatrix} = 
\begin{bmatrix}
\cos \theta & -v \sin \theta \\
\sin \theta & v \cos \theta
\end{bmatrix}
\begin{bmatrix}
\dot{v} \\
\omega
\end{bmatrix}.
\] (18)

Inverting the transformation (18) we obtain the input that must be applied to the robot for it to behave as a linear system

\[
\begin{bmatrix}
\dot{v} \\
\omega
\end{bmatrix} = 
\begin{bmatrix}
\cos \theta & \sin \theta \\
-\frac{1}{v} \sin \theta & \frac{1}{v} \cos \theta
\end{bmatrix}
\begin{bmatrix}
u_1 \\
u_2
\end{bmatrix}.
\] (19)

Using (18) and (19), system (14) is equivalent to (17), i.e., \( x = w_1, \dot{x} = w_1, y = w_2, \dot{y} = w_2 \).

Care must be taken when using (19), since it is not defined for \( v = 0 \). This is known to be a structural problem of dynamic extensions of wheeled mobile robots [28]. It should be taken into account when designing control laws for the equivalent linear system.

VI. SIMULATION EXAMPLES

In this section, we consider six agents interacting under the communication topology of Fig. 2. The \( C \) matrix corresponding to this topology has the eigenvalues \( \{1, 0.46, -0.5, -0.64, -0.16 + j0.21\} \) which create four second-order factors and one fourth-order factor, according to (3). The number of agents was selected because this is the number of actual experimental platforms that are available for a follow-on experimental study. The communication topology was selected such as its \( C \) matrix presented both real and complex eigenvalues without repetitions. For a larger number of agents the analysis methodology remains the same, but the computational load may increase, as larger communication topologies may create more diverse sets of eigenvalues.

By deploying the CTCR paradigm into the factors, considering control gains \( P = 1 \) and \( D = 0.5 \), we obtain the stability chart of Fig. 3 in the domain of the time delays. Each one of the factors creates some kernel, depicted in red, and offspring curves, depicted in blue. At each point on those curves the system has a pair of imaginary characteristic roots. By increasing either of the delays at one of those points the number of unstable characteristic roots changes. The shaded region in that plot corresponds to combinations of delays \( (\tau_1, \tau_2) \) for which the number of unstable roots is zero. That is, this is the stable operating region. For any delay combination within the gray area, no factor of the characteristic equation of the system has an unstable root.

We wish to highlight that this chart is exact and exhaustive: no delay combination outside of the gray area results in stable operation, and those in the boundary of the gray area create margin-
ally stable operation, with the agents oscillating around a fixed position.

To test the formation control algorithm, we propose the formation structure of Fig. 4. For this configuration, the desired values are $[0.71, 1.83, 1.41, -2.08, -0.81, 5.74]$ in the horizontal direction and $[-1.12, -0.44, 1.41, 1.50, -1.40, 3.72]$ in the vertical direction. In the simulation, the agents start from random positions and orientations, and all have very small initial speeds to avoid the mentioned division by zero in (19) at $v = 0$. Using the delay combination $(\tau_1, \tau_2) = (0.3 \text{ s}, 0.2 \text{ s})$, which corresponds to point a in Fig. 3, the agents settle into the formation after about 30 s. Fig. 5 shows the positions in $x$ and $y$ direction as a function of time, whereas Fig. 6 shows how the agents reach the desired formation as final configuration. The short black lines represent the instantaneous heading of the agents at the moment of the snapshot.

An increment of the delays generally worsens the performance of the system. When both delays are increased to $(\tau_1, \tau_2) = (2 \text{ s}, 1 \text{ s})$, corresponding to point b in Fig. 3, the agents take more than 100 s to reach the final state, as can be observed in Fig. 7.
Increasing the delays, however, is not always bad. An interesting, and rather counterintuitive fact that Fig. 3 shows is that, in some regions of the delay space, stability can be recovered by increasing one, or both, delays. For example, the delay selection \((\tau_1, \tau_2) = (1 \text{ s}, 2 \text{ s})\), marked by point e in Fig. 3, is in the unstable region. By either increasing \(\tau_2\) to 5.5 s (point d) or \(\tau_1\) to 3.5 s (point e), the stability is recovered. This is a very important result, since it provides an extra degree of flexibility to the control designer. Reducing the delays is normally not possible without changing the hardware, but increasing them is simple.

This idea of delay scheduling was first discussed in [30] and [31]. It basically consists on a skillful increase of the delays to maintain stability or improve the performance of a system. In the case under study, we would be interested in the delay combination that produces the fastest convergence to the formation. For this, we need to calculate the dominant root of the system, and its associated time constant, for every delay combination within the stable operating zone of Fig. 3. This can be done only with numerical procedures, like the QPmR algorithm [32]. The stability analysis using the CTCR paradigm is a necessary and crucial first step for delay scheduling, as it is the only methodology that exhaustively declares the stable regions of the system, and therefore provides a finite area in which to search for the smallest time constant.

For the system under study, Fig. 8 shows how the real part of the dominant root of the system changes as function of the delays within a domain that includes part of the stability region. The boundary depicted in Fig. 3 can also be seen in Fig. 8 as the intersection of the surface with the plane defined by \(\Re(s) = 0\).

The deepest valleys of the surface in Fig. 8 correspond to the smallest time constants, and therefore to the fastest formation convergence rates. The plot shows that the best delay combinations are those in the vicinity of the origin, and in a small region for \(\tau_1 \in (0, 1) \text{ s}\) and \(\tau_2 \in \{4, 6\} \text{ s}\). This fact is better observed in Fig. 9, which presents a contour map of the surface in Fig. 8.

With the information contained in Figs. 8 and 9, the designer can properly tune the delays to achieve a desired response time for the system. This is the core idea of the delay scheduling.

This paper considered the problem of cooperative formation control for nonholonomic agents in the presence of time delays. We deployed a formation control algorithm based on a linear consensus protocol. We consider two time delays: one in the exchange of position information and the other in the exchange of velocity information. Since the control law was originally devised for linear, holonomic systems, a feedback linearization technique is used to adapt the control logic to the class of agents under study.

The stability of the formation control depends on the stability of the consensus protocol. In order to find the delay combinations that lead to stable operation, the CTCR paradigm was used. Before using CTCR, a state decomposition that leads to a factorization of the characteristic equation of the system is performed to simplify the problem. With this combination of factorization plus CTCR, it is possible to obtain, in a very efficient manner, an exact and exhaustive stability map with respect to the delays. A simulation shows that the convergence to formation is guaranteed in the regions declared by the proposed methodology.

We also discussed briefly the concept of delay scheduling and how the delays can be used to tune the time the agents take to reach the desired formation.

The codes used to create the examples are available at https://db.tt/JhOOzQCF or can be requested via email to the authors.

Future work includes the practical implementation of these algorithms in a team of mobile robots. For this, the Process Control and Automation Laboratory of the Universidad Santo Tomás is working on the development of an experimental platform based on the iRobot Create development kit [26].

REFERENCES


